

# Dealing with Ambiguity

- ▶ Consider possible parses but weighted by probability
- ▶ Return likeliest parse
- ▶ Return likeliest parse along with a probability

# PCFG: Probabilistic Context-Free Grammar

- ▶ Components of PCFG:  $G = \langle N, \Sigma, R, S \rangle$ 
  - ▶  $\Sigma$ , an alphabet or set of *terminal* symbols
  - ▶  $N$ , a set of *nonterminal* symbols,  $N \cap \Sigma = \emptyset$
  - ▶  $S \in N$ , a *start* symbol (distinguished nonterminal)
  - ▶  $R$ , a set of *rules* or *productions* of the form

$$A \longrightarrow \beta[p]$$

- ▶  $A \in N$  is a single nonterminal and  $\beta \in (\Sigma \cup N)^*$  is a finite string of terminals and nonterminals
- ▶  $p = P(A \longrightarrow \beta | A)$  is the probability of expanding  $A$  to  $\beta$

$$\sum_{\beta} P(A \longrightarrow \beta | A) = 1$$

- ▶ Consistency:
  - ▶ Probability of a sentence is nonzero if and only if it is in the language
  - ▶ Sum of probabilities of sentences in the language is 1

# Languages from Grammars

- ▶ Simple CFG: Nominal is the start symbol

Nominal  $\rightarrow$  Nominal Noun

Nominal  $\rightarrow$  Noun

Noun  $\rightarrow$  olive

Noun  $\rightarrow$  jar

- ▶ Simpler CFG: Nominal is the start symbol

Nominal  $\rightarrow$  Nominal Noun

Noun  $\rightarrow$  olive

Noun  $\rightarrow$  jar

- ▶ Simple PCFG: Nominal is the start symbol

Nominal  $\rightarrow$  Nominal Noun  $[\frac{2}{3}]$

Nominal  $\rightarrow$  Noun  $[\frac{1}{3}]$

Noun  $\rightarrow$  jar  $[1]$

# Consistent PCFG

Probability of the language is 1

- ▶ Consider the same simple PCFG as before
  - Nominal  $\rightarrow$  Nominal Noun  $[\frac{2}{3}]$
  - Nominal  $\rightarrow$  Noun  $[\frac{1}{3}]$
  - Noun  $\rightarrow$  jar  $[1]$
- ▶ Write out all parse trees for  $\text{jar}^k$
- ▶ Probability of  $\text{jar}^k$  is sum of probabilities for its parse trees
- ▶ Sum up the probabilities for the entire language

# Inconsistent PCFG

Probability of generating the language is not 1

- ▶ Consider a modified PCFG: Nominal is the start symbol  
Nominal  $\rightarrow$  Nominal Nominal  $[\frac{2}{3}]$   
Nominal  $\rightarrow$  jar  $[\frac{1}{3}]$
- ▶ Write out all parse trees for  $\text{jar}^k$
- ▶ Probability of  $\text{jar}^k$  is sum of probabilities for its parse trees
- ▶ Sum up the probabilities for the entire language

The argument gets cumbersome

# PCFG: Markovian Argument

- ▶ Consider how a derivation proceeds
    - ▶ One production increases the count of nonterminals by one
    - ▶ One production decreases the count of nonterminals by one
    - ▶ We start with one nonterminal (the start symbol)
    - ▶ Any derivation that ends in zero nonterminals yields a string in the language
  - ▶  $L(n+1)$  (left move): probability of starting from  $n+1$  nonterminals and arriving at a state with  $n$  nonterminals
- The probability of generating a string in this language is  $L(1)$
- ▶  $L(0)$  is never used and could be left undefined or set to zero
  - ▶ PCFGs respect the Markov assumption: any nonterminal has an equal chance of being expanded regardless of history
  - ▶ Therefore,  $L(n+1)$  is a constant,  $L$

# Inconsistent PCFG: Markovian Derivation

- ▶ Probabilities of stepping right  $q$  and left  $1 - q$
- ▶  $L$  (probability of eventually moving one left) equals
  - ▶ Stepping one left immediately plus
  - ▶ Stepping one right followed by two paths moving one step left each

$$L = 1 - q + qL^2$$

- ▶ Solve  $qL^2 - L + 1 - q = 0$
- ▶  $L = \frac{1 \pm \sqrt{1 - 4q(1 - q)}}{2q}$
- ▶  $\sqrt{1 - 4q(1 - q)} = (2q - 1)$
- ▶ Therefore,  $L$  has two solutions, of which the *minimum* is appropriate
  - ▶ Trivial solution:  $L = \frac{1 - (1 - 2q)}{2q} = 1$
  - ▶ Left-right odds:  $L = \frac{1 - (2q - 1)}{2q} = \frac{1 - q}{q}$
- ▶ For our example,  $L = \min(1, \frac{1}{3}) = \frac{1}{3} \neq 1$ —indicating inconsistency
- ▶ If we reverse the probabilities, then  $\min(1, 2) = 1$

# Probability of a Parse Tree

- ▶ Tree  $T$  obtained from sentence  $W$ , i.e.,  $T$  yields  $W$

$$P(T, W) = P(T)P(W|T)$$

$$P(T, W) = P(T) \text{ since } P(W|T) = 1$$

- ▶ Obtaining  $T$  via  $n$  expansions  $A_i \rightarrow \beta_i$  and  $S = A_1$  is the start symbol

$$P(T, W) = \prod_{i=1}^n P(\beta_i | A_i)$$

- ▶ Best tree for  $W$

$$\hat{T}(W) = \operatorname{argmax}_{T \text{ yields } W} P(T|W) = \operatorname{argmax}_{T \text{ yields } W} \frac{P(T, W)}{P(W)}$$

- ▶ Since  $P(T, W) = P(T)$  and  $P(W)$  is constant ( $W$  being fixed)

$$\hat{T}(W) = \operatorname{argmax}_{T \text{ yields } W} P(T)$$



# Probabilistic CKY Parsing

- ▶ Like CKY, as discussed earlier, except that
  - ▶ Each cell contains not a set of, but a probability distribution over, nonterminals
- ▶ Specifying probabilities for Chomsky Normal Form
  - ▶ Consider each transformation used in the normalization
- ▶ Supply the probabilities below
  - ▶ Replace  $A \rightarrow \alpha B \gamma [p]$  and  $B \rightarrow \beta [q]$  by  $A \rightarrow \alpha \beta \gamma [?]$
  - ▶ Replace  $A \rightarrow BC \gamma [p]$  by  $A \rightarrow BX [?]$  and  $X \rightarrow C \gamma [?]$
- ▶ Store a probability distribution over nonterminals in each cell
- ▶ Return likeliest parse

# Learning PCFG Probabilities

- ▶ Simplest estimator: Assume a treebank
- ▶ Estimate the probability of  $A \rightarrow \beta$  as

$$P(A \rightarrow \beta | A) = \frac{\text{Count}(A \rightarrow \beta)}{\sum_{\gamma} \text{Count}(A \rightarrow \gamma)} = \frac{\text{Count}(A \rightarrow \beta)}{\text{Count}(A)}$$

- ▶ Without a treebank but with a corpus
- ▶ Assume a traditional parser
- ▶ Initialize all rule probabilities as equal
- ▶ Iteratively
  - ▶ Parse each sentence in the corpus
  - ▶ Credit each rule  $A \rightarrow \beta_i$  by the counts weighted by the probabilities of the rules leading to that nonterminal,  $A$
  - ▶ Revise the probability estimates
- ▶ More properly described as an expectation maximization algorithm

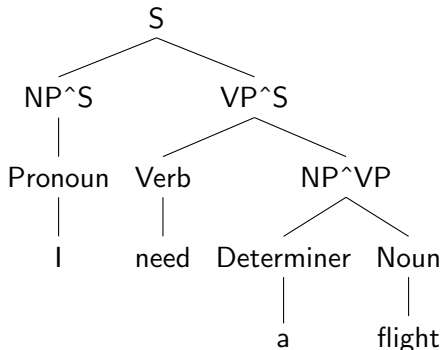
# Shortcomings of PCFGs

PCFGs break ties between rules in a fixed manner

- ▶ Naïve context-free assumption regarding probabilities
  - ▶  $NP \xrightarrow{*}$  Pronoun much likelier for a Subject NP than an object NP
  - ▶ PCFGs (and CFGs) disregard the path on which the NP was produced
- ▶ Lack of lexical dependence
  - ▶  $VP \rightarrow VBD\ NP\ NP$  is likelier for a ditransitive verb
- ▶ Consider prepositional phrase attachment
  - ▶ Either: prefer PP attached to VP (“dumped sacks into a bin”)
    - ▶  $VP \rightarrow VBD\ NP\ PP$
  - ▶ Or: prefer PP attached to NP (“caught tons of herring”)
    - ▶  $VP \rightarrow VBD\ NP$
    - ▶  $NP \rightarrow NP\ PP$
- ▶ Coordination ambiguities: each parse gets the same probability because all parses use the same rules

## Split Nonterminals to Refine a PCFG

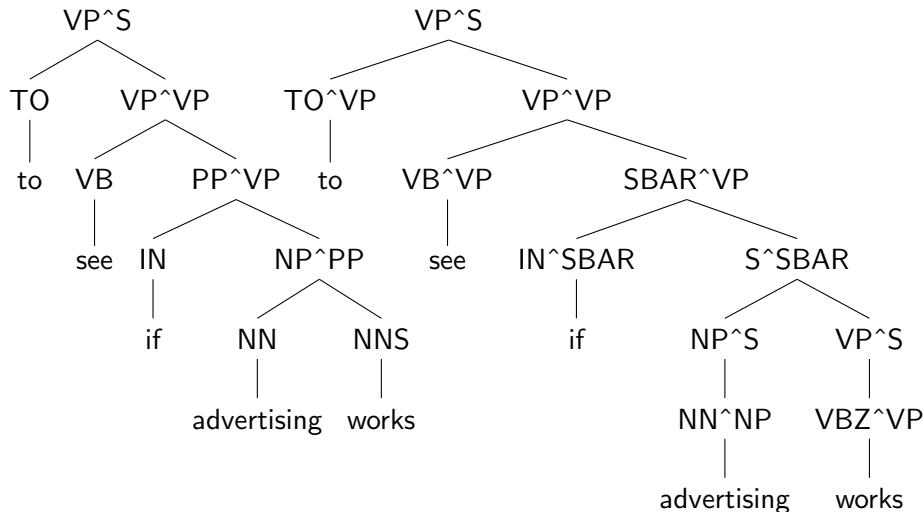
- ▶ Split nonterminals for syntactic roles, e.g.,  $NP_{\text{subject}}$  versus  $NP_{\text{object}}$ 
  - ▶ Then learn different probabilities for their productions
- ▶ Capture part of path by a parent annotation
  - ▶ Annotating only the phrasal nonterminals ( $NP^S$  versus  $NP^{VP}$ )



- ▶ Likewise, split *preterminals*, i.e., nonterminals that yield terminals
  - ▶ Adverbs depend on where they occur:  $RB^{AdvP}$  (also, now),  $RB^{VP}$  (not),  $RB^{NP}$  (only, just)

# Example of Preterminals with Sentential Complements

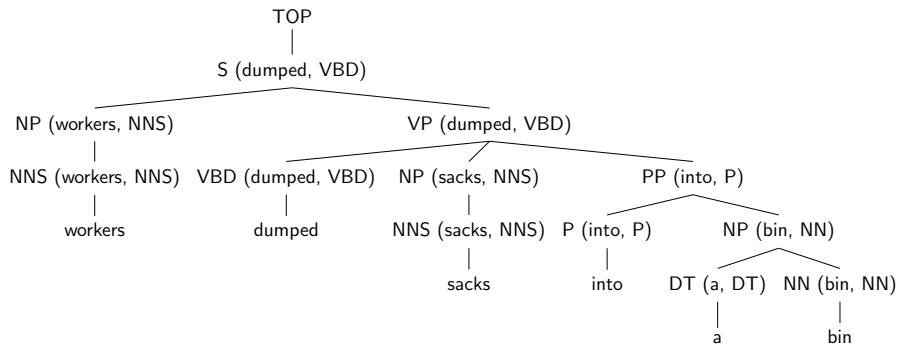
Klein and Manning: Left parse is wrong



IN includes preps, complementizers (that), subord conjs (if, as)

# Lexicalized Parse Tree

Variant of previous such tree with parts of speech inserted




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TOP → S(dumped, VBD)

S(dumped, VBD) → NP(workers, NNS) VP(dumped, VBD)

VP(dumped, VBD) → VBD(dumped, VBD) NP(sacks, NNS) PP(into, P)

...

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VBD(dumped, VBD) → dumped

...

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# Estimating the Probabilities

- ▶ In general, we estimate the probability of  $A \rightarrow \beta$  as

$$P(A \rightarrow \beta | A) = \frac{\text{Count}(A \rightarrow \beta)}{\sum_{\gamma} \text{Count}(A \rightarrow \gamma)} = \frac{\text{Count}(A \rightarrow \beta)}{\text{Count}(A)}$$

- ▶ But the new productions are highly specific
- ▶ Collins Model 1 makes independence assumptions
  - ▶ Treat  $\beta$  as  $\beta_1 \dots \beta_H \dots \beta_n$ :  $\beta_H$  is the head and  $\beta_1 = \beta_n = \text{STOP}$
  - ▶ Generate the head
  - ▶ Generate its premodifiers until getting to STOP
  - ▶ Generate its post-modifiers until getting to STOP
  - ▶ Apply Naïve Bayes

$$P(A \rightarrow \beta) = P(A \rightarrow \beta_H) \times P(\beta_1 \dots \beta_{H-1} | \beta_H) \times P(\beta_{H+1} \dots \beta_n | \beta_H)$$

$$\approx P(A \rightarrow \beta_H) \times \prod_{k=1}^{H-1} P(\beta_k | \beta_H) \times \prod_{k=H+1}^n P(\beta_k | \beta_H)$$

- ▶ Estimate each probability from smaller amounts of data

# Labeled Recall and Precision to Evaluate Parsers

- ▶ Like recall and precision but
  - ▶ Based on counting correct constituents identified
  - ▶ Correctness with respect to a ground truth *reference* parse tree
- ▶ Recall
  - ▶ How many of the correct constituents are discovered
- ▶ Precision
  - ▶ How many of the constituents discovered are correct



# Cross Brackets

A metric specific to comparing parse trees

- ▶ A measure of error
- ▶ The number of constituents for which
  - ▶ The reference parse has a bracketing ((A B) C)
  - ▶ The hypothesis parse has a bracketing (A (B C))
- ▶ On the Wall Street Journal treebank, modern parsers yield
  - ▶ Recall 90%
  - ▶ Precision 90%
  - ▶ Cross-bracketing 1%
- ▶ Extended metrics for comparing parsers using different grammars

# Human Parsing

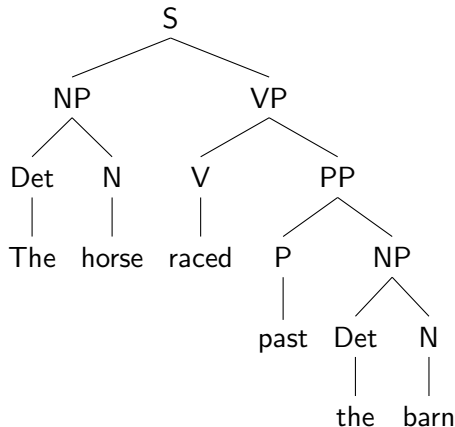
## Psycholinguistics

- ▶ Studies of human processing ease
  - ▶ Delay in reading
  - ▶ Eye gaze fixation (dwell) time
- ▶ Garden-path sentences
  - ▶ Prefix (initial portion) is ambiguous
  - ▶ That is, temporarily ambiguous while reading
  - ▶ A higher preferred parse of the prefix doesn't lead to a parse of the entire sentence

# The Horse Raced Past the Barn Fell: Problematic

A complete sentence followed by an extra verb

The first part gets a likely parse that offers no clear attachment for the final verb



# The Horse Raced Past the Barn Fell: Correct

Raced is part of a reduced relative clause modifying “The horse”

